Llull and Copeland Voting Broadly Resist Bribery and Control

Piotr Faliszewski
Department of Computer Science
University of Rochester
Rochester, NY 14627

Edith Hemaspaandra
Department of Computer Science
Rochester Institute of Technology
Rochester, NY 14623

Lane A. Hemaspaandra
Department of Computer Science
University of Rochester
Rochester, NY 14627

Jörg Rothe†
Institut für Informatik
Heinrich-Heine-Universität Düsseldorf
40225 Düsseldorf, Germany

February 19, 2007

Abstract

Control of elections refers to attempts by an agent to, via such actions as addition/deletion/partition of candidates or voters, ensure that a given candidate wins [BTT92]. An election system in which such an agent’s computational task is NP-hard is said to be resistant to the given type of control. The only election systems known to be resistant to all the standard control types are highly artificial election systems created by hybridization [HHR07]. In this paper, we prove that an election system developed by the 13th century mystic Ramon Llull and the well-studied Copeland election system are both resistant to all the standard types of (constructive) electoral control other than one variant of addition of candidates. This is the most comprehensive resistance to control yet achieved by any natural election system. In addition, we show that Llull and Copeland voting are very broadly resistant to bribery attacks, and we integrate the potential irrationality of voter preferences into many of our results.

1 Introduction

Elections have played an important role in human societies for thousands of years. For example, elections were of central importance in the democracy of ancient Athens. There, citizens typically could only agree (vote yes) or disagree (vote no) with the speaker, and

*Supported in part by DFG grant RO-1202/9-3, NSF grants CCR-0311021 and CCF-0426761, the Alexander von Humboldt Foundation’s TransCoop program, and a Friedrich Wilhelm Bessel Research Award. A more complete, with-full-proofs version of this report is in preparation and will appear at the URCS TR web site in March or April, 2007.

†Work done in part while visiting the University of Rochester.
simple majority-rule was in effect. Mathematical study of elections, give or take a few discussions by the ancient Greeks, was until recently thought to have been initiated only a few of hundred years ago, namely, in the breakthrough work of Borda and Condorcet—later in part reinvented by Dodgson. One of the most interesting results of this early work is Condorcet’s observation that if one conducts elections with more than two alternatives then even if all voters have rational preferences, the society as a whole might behave irrationally (i.e., for any alternative, a majority of the voters would prefer some other option to win). Based on his observations, Condorcet suggested that if there exists a candidate $c$ such that $c$ defeats any other candidate in a head-to-head contest then that candidate should win the election. Such a candidate is called a Condorcet winner. Clearly, there can be at most one Condorcet winner in any election and there might be none.

This understanding of history has been shattered during the past few decades, as it has been rediscovered that the study of elections was in fact considered deeply as early as the thirteenth century (see Hägèle and Pukelsheim [HP01] and the citations therein regarding Ramon Llull and the fifteenth century figure Cusanus, esp. the citations that there are numbered 3, 5, and 24–27). Ramon Llull (b. 1232, d. 1315), a Catalan mystic, missionary, and philosopher developed an election system that (a) has an efficient winner-determination procedure and that (b) elects a Condorcet winner whenever one exists and otherwise elects candidates that are, in some sense, closest to being Condorcet winners. Llull’s motivation for developing an election system was to obtain a method of choosing the abbesses, abbots, bishops, and perhaps even the pope. His ideas never gained public acceptance in medieval Europe and were long forgotten.

Llull’s election system is remarkably close to that of Copeland. For each pair of distinct candidates, Llull asks all the voters which among the two they prefer. After such a head-to-head contest is conducted, the candidate(s) that got at least half of the votes get one point. (Note that if the candidates tie then they both get a point; in Copeland’s system in such a case neither one gets a point.\footnote{Page 23 of Hägèle and Pukelsheim [HP01] indicates in a way we find deeply convincing (namely by a direct quote of Llull’s in-this-case-very-clear words from his Artifitium Electionis Personarum—which was rediscovered by those authors in the year 2000) that at least one of Llull’s election systems was defined in this way, and so in this paper we refer to both-candidates-score-a-point-on-a-tie as Llull voting. To avoid any confusion, we will retain the standard usage of referring to the case when head-to-head ties score an overall point for neither candidate as Copeland elections, though we mention that Llull, who proposed multiple systems and was not always clear, has also been credited by McLean and Lorrey [ML06], drawing on different words in the same manuscript, for inventing what is now called Copeland voting.} This minor difference can make the dynamics of Llull’s system quite different from those of Copeland’s.) In the end, the candidates with the largest number of points are declared winners. (Llull in some settings required the candidate and voter sets to be identical, and had an elaborate two-stage tie-breaking rule ending in randomization, but here we take simply his underlying system, without tie-breaks and with the voter and candidate sets unrestricted. That is, we in these matters cast his system into the modern idiom for election systems.)

It is interesting to note that Llull in fact allowed voters to have irrational preferences. Given three candidates, $c$, $d$, and $e$, it was perfectly acceptable for a voter to prefer $c$ to $d$, ...
to $e$, and from $e$ to $c$. On the other hand, in modern studies of voting and election systems each voter’s preferences are most typically modeled as a strict linear order over all candidates. Yet irrationality is a very tempting and natural concept. Consider Bob, who likes to eat out but is often in a hurry. Bob prefers diners to fast food because he is willing to wait a little more to get better food. Also, given a choice between a fancy restaurant and a diner he prefers the restaurant; again because he is willing to wait somewhat longer to get better quality. However, given the choice between a fast-food place and a fancy restaurant Bob might reason that he is not willing to wait so long for the fancy restaurant and will choose fast food instead. Thus regarding catering options, Bob’s preferences are irrational in our sense. Similar irrationalities easily come up when voters make their choices based on multiple criteria—a very natural scenario both among humans and software agents.

The goal of this paper is to study Llull’s and Copeland’s election systems from the point of view of computational social choice theory, in the setting where voters are rational and in the setting where the voters are allowed to have irrational preferences. (Note: When we henceforth say “irrational voters,” we mean that they may have irrational preferences, not that they each must.)

In general it is impossible to design a perfect election system. In the 1950s Arrow famously showed that there is no social choice system that satisfies a certain small set of requirements, and later Gibbard, Satterthwaite, and Duggan and Schwartz showed that any natural election system can be manipulated by strategic voting, i.e., by a voter who reveals different preferences than her true ones in order to affect an election’s result in her favor. Also, no natural election system has yet been shown to be resistant to all types of control via procedural changes.

These obstacles are very depressing, but the field of computational social choice theory grew in part from the realization that computational complexity provides a tool to partially circumvent them. In particular, around 1990 Bartholdi, Tovey, and Trick [BTT89, BTT92], and Bartholdi and Orlin [BO91] brilliantly observed that while we might not be able to make manipulation (i.e., strategic voting) and control of elections impossible, we can at least try to make such manipulation and control so difficult computationally that neither voters nor election organizers will attempt it. For example, if there is a way for a committee’s chair to set up an election within her committee in such a way that her favorite option is guaranteed to win but the chair’s computational task would take a million years, then for all practical purposes we may assume that the chair is prevented from finding such a set-up.

So since the seminal work of Bartholdi, Orlin, Tovey, and Trick, a large body of research has naturally been dedicated to the study of computational properties of the election systems. Some topics that received much attention are the complexity of manipulation of elections [CS02a, CS03, CS02b, CS06, CLS03, EL05, HH07, PR06, PRZ07] and the control of elections via procedural changes [HHR05, HHR07, PRZ07]. Recently, Faliszewski, Hemaspaandra, and Hemaspaandra [FHH06] studied the complexity of bribery in elections. Bribery shares some features of manipulation and some of control. In particular, the briber picks the voters he or she wants to affect (as in voter control problems) and asks them to vote as he or she wishes (as in manipulation).
In this paper we study Llull and Copeland elections with respect to the difficulty of bribery and procedural control. We believe that the computational complexity of procedural control within elections is an important topic in multiagent systems. In multiagent systems, elections can be used as a tool to solve many practical problems. As just a few of examples we mention the work of Ephrati and Rosenschein [ER93] where elections are used for the purposes of planning, and the work of Dwork et al. [DKNS01] where elections are used to aggregate results from multiple web-search engines. In a multiagent setting we might have hundreds of elections happening every minute and we cannot hope to carefully check if in each case the party that organized the elections did not attempt some procedural changes to skew the results. However, if it is computationally hard to effectuate such procedural changes then we can hope it is impossible for the organizers to undertake them.

A standard technique for showing that a particular elections-related problem (e.g., the problem of deciding whether the chair can make her favorite candidate \( p \) a winner by influencing at most \( k \) voters not to cast their votes) is computationally difficult is to show that it is NP-hard. This approach is taken in all of the papers on computational social choice theory cited above, and it is the approach that we take in this paper. One of the justifications for using NP-hardness as a measure of difficulty is that in multiagent settings manipulation and control of elections is conducted by computationally bounded software agents that neither have human intuition nor the computational ability to solve NP-hard problems.

Our bribery results and some of our control results are proven by a method we introduce of controlling the relative performances of certain voters in such a way that, if one sets up other restrictions appropriately, the legal possibilities for control/change actions are sharply constrained. We call our approach “the UV technique,” since it is based on dummy candidates \( U \) and \( V \). The proof of Lemma 4 is a particular case of the method. We feel that the UV technique will be useful even beyond the bounds of this paper, and we now provide an example showing that. In particular, the authors have recently noticed that the proof of Bartholdi, Tovey, and Trick [BTT92, Theorem 12] of an important result of theirs related to this paper (namely, that Condorcet voting is resistant to constructive control by deleting voters) is invalid. The invalidity is due to the proof centrally using nonstrict voters, in violation of Bartholdi, Tovey, and Trick’s [BTT92] (and our) model, and the invalidity seems very hard to fix with the proof approach taken there. However, using the UV technique one indeed can easily obtain a correct proof, and we have done so (our in-preparation full version will present the proof). Thus the UV technique this paper introduces has applications even beyond those in this paper, and in particular, Theorem 12 of [BTT92] holds. We noticed that Theorem 14 of the same paper also has a similar flaw and we have validly proven it also, again using the UV technique. The UV technique in some sense provides a more uniform approach for (certain) voter-related election problems.

This paper is organized as follows. In the Preliminaries section we formalize the notions of Llull and Copeland elections and introduce our notation. In the Bribery section we will show that Llull and Copeland are perfect from the point of view of resistance to bribery, both in the case of rational voters and in the case of irrational voters. On the other hand,
we show that if one changes the bribery model to allow “micro-bribes” of voters, their
resistance in the irrational case is pierced by the existence of a polynomial-time algorithm.
In the Control section we present our results on the procedural control of Llull and Copeland
elections. What this section shows is, in effect, that both Llull and Copeland have resistance
to more types of (constructive) control than has been shown for any other known natural
system.

2 Preliminaries

An election consists of a candidate set $C = \{c_1, \ldots, c_n\}$, a collection $V$ of voters, and a rule
that specifies the winners. As is standard, each voter is represented (individually, except
later when we discuss succinct inputs) via her preferences. We consider two ways in which
voters can express their preferences. In the rational case, each voter’s preferences are a strict
linear order over the set $C$, e.g., each voter has a preference list $c_{i_1} > c_{i_2} > \cdots > c_{i_n}$, with
the $i_j$’s a permutation. In the irrational case, each voter’s preferences are represented as a
table that for every unordered pair of distinct candidates $c_i$ and $c_j$ in $C$ indicates whether
the voter prefers $c_i$ to $c_j$ (i.e., $c_i > c_j$) or $c_j$ to $c_i$ (i.e., $c_j > c_i$).

Some well-known election rules for the case of rational voters include plurality, Borda
count, and Condorcet. Plurality elects the candidate(s) that are ranked highest by the
largest number of voters. Borda count elects the candidate(s) that receive the most points,
where each voter $v_i$ gives each candidate $c_j$ as many points as the number of candidates
$c_j$ defeats on $v_i$’s preference list. A Condorcet winner is a candidate $c_i$ such that for every
other candidate $c_j$ it holds that $c_i$ is preferred to $c_j$ by a strict majority of the voters. Note
that all of these election systems can easily be adapted to work with irrational voters.

The central focus of this paper is on the tournament-based systems of Llull and
Copeland, for which definitions were given during the Introduction.

Let us now describe the computational problems that we study in this paper. Our
problems come in two flavors: constructive and destructive. In the constructive version the
goal is to test whether, via either bribery or control, it is possible to make a designated
candidate a winner of the election. In the destructive case we seek to guarantee that a
despised candidate is not a winner of the election.

The (constructive) bribery problem for the Copeland system with rational voters, which
we call Copeland-bribery, is defined as:

**Given:** A set $C$ of candidates, a collection $V$ of voters specified via their preference lists,
a distinguished candidate $p$, and a nonnegative integer $k$.

**Question:** Is it possible to make $p$ a winner of the Copeland election by modifying the
preference lists of at most $k$ voters?

The same problem for Copeland with irrational voters is called
Copeland_{Irrational}-bribery, and the corresponding problems for Llull are called Llull-bribery
for the rational case and Llull_{Irrational}-bribery for the irrational case.
Bribery problems seek to change the outcome of elections via modifying the preferences of some of the voters. On the other hand, control problems seek to change the outcome of elections by modifying their structure, e.g., via adding candidates (AC), deleting candidates (DC), partitioning candidates (PC), partitioning candidates with run-off (RPC), adding voters (AV), deleting voters (DV), and partitioning voters (PV). The name of a control problem is formed by concatenating the name of the election system with CC or DC, for constructive control and destructive control respectively, and the acronym of the control type that we have in mind. For the case of partitioning candidates and voters we also distinguish two subcases, ties-eliminate (TE) and ties-promote (TP). In the ties-eliminate subcase, if a subelection is won by more than one candidate then that subelection doesn’t promote any of its candidates to the next round. In the ties-promote subcase, all the winners within a given subelection are promoted to the next round. We below give a formal description of just a few of the above control types, and via their names have above described the rest of the control types informally; interested readers can find the missing control-type definitions in Hemaspaandra, Hemaspaandra, and Rothe [HHR05].

Copeland destructive control via deleting voters for the case of rational voters, Copeland-DCDV, is defined as:

**Given:** A set $C$ of candidates and a collection $V$ of voters represented via preference lists over $C$, a distinguished candidate $c$, and a nonnegative integer $k$.

**Question:** Is it possible to by deleting at most $k$ voters ensure that $c$ is not a winner of the resulting Copeland election?

Similarly, LlullIrrational-CCAC, the problem of constructive control of Llull elections with irrational voters via adding candidates is defined as:

**Given:** Disjoint sets $C$ and $D$ of candidates, a collection $V$ of voters specified via their (possibly irrational) preference tables over the candidates in the set $C \cup D$, and a distinguished candidate $p$.

**Question:** Is it possible to choose a subset $E$ of $D$ such that $p$ is a winner of the Llull election with voters $V$ and candidates $C \cup E$?

The above definition of LlullIrrational-DCAC is based on that introduced by Bartholdi, Tovey, and Trick [BTT92]. In contrast with the other control problems involving adding or deleting candidates or voters, in the adding candidates problem Bartholdi, Tovey, and Trick did not introduce a nonnegative integer $k$ that bounds the number of candidates (from the set $D$) the chair is allowed to add. We feel this is somewhat inconsistent and thus we define a with-change-parameter version of the control by adding candidates problems: In the $E$-k-CCAC we ask if it is possible to make the distinguished candidate $p$ a winner of election $E$ by adding at most $k$ candidates from the spoiler candidate set $D$. We define the destructive version, $E$-k-DCAC, analogously.
Note that our problems—both those regarding bribery and those regarding control—speak of nonunique winners. Nonetheless, we have proven all our control results both in the case of nonunique and (to be able to fairly compare them with existing control results, which except for the interesting “multiwinner” model of Procaccia, Rosenschein, and Zohar [PRZ07] are for the unique winner case) unique winners.

Not all election systems can be affected by each control type. For example, if a candidate \( c \) is a Condorcet winner then it is impossible to prevent her from being a Condorcet winner by deleting other candidates. However, for the case of Llull and Copeland systems it is easy to see that for each standard type of control there is a scenario where the outcome of the election can be changed via conducting the control action.

We say that an election system (Llull, Copeland, etc.) is resistant to a particular attack (be it a type of control or of bribery) if the appropriate computational problem is NP-hard and (only for the control case) there is a scenario where this type of attack can change the winners of the election in the appropriate constructive/destructive way (note: for Llull and Copeland the second clause, needed for some systems with complex evaluation problems, is superfluous). On the other hand, if the computational problem is in P and there is a scenario where this type of attack can change the winners of the election, then we say the system is vulnerable to this attack. For all resistance claims in this paper, NP-membership is clear, and so NP-completeness holds. The notions of resistance and vulnerability (and of immunity and susceptibility) for control problems in election systems were introduced by Bartholdi, Tovey, and Trick [BTT92].

3 Bribery

Our main result regarding bribery is that the Llull and Copeland systems are resistant to bribery regardless of voters’ rationality and our mode of operation (constructive versus destructive).

**Theorem 1** The Llull and Copeland systems are resistant to both constructive and destructive bribery in both the rational-voters case and the irrational-voters case.

Theorem 1 follows by an application of our UV technique. The application of the UV technique here is very similar to that in Lemma 4, whose proof we include in the next section.

There is a natural yet different and more micro-scale way of defining bribery for the case of irrational voters. The way our bribery problems are defined, we can choose up to \( k \) voters and modify their preferences freely. However, what if, for the case of irrational voters, instead of modifying the whole preference table of some selected voters we now pay unit cost each time we flip a single entry in a voter’s table? Call this version of the problem bribery’. Theorem 1 notwithstanding, the bribery’ problems are easy.

**Theorem 2** The Llull and Copeland systems with irrational voters are vulnerable to constructive and destructive bribery’.
The destructive cases follow via greedy algorithms. On the other hand, our algorithms for the constructive cases are rather involved. Very briefly put, the main idea is that we model Llull/Copeland elections via a network flow problem, where the units that flow in the network are Llull/Copeland points. This allows us to use min-cost flow problem algorithms to model even very complicated interactions among candidates in the election.

4 Control

In this section we focus on control in Llull and Copeland elections, and we compare our results to plurality elections.

Table 1 presents our results on the resistance and vulnerability of Llull and Copeland elections to procedural control.

**Theorem 3** Llull and Copeland elections are resistant and vulnerable to control types as indicated in Table 1. The same results, regarding Lull and Copeland, hold for the case of irrational voters.

<table>
<thead>
<tr>
<th>Control type</th>
<th>Llull</th>
<th>Copeland</th>
<th>Plurality</th>
</tr>
</thead>
<tbody>
<tr>
<td>AC</td>
<td>V</td>
<td>V</td>
<td>V</td>
</tr>
<tr>
<td>k-AC</td>
<td>R</td>
<td>V</td>
<td>R</td>
</tr>
<tr>
<td>DC</td>
<td>R</td>
<td>V</td>
<td>R</td>
</tr>
<tr>
<td>RPC-TP</td>
<td>R</td>
<td>V</td>
<td>R</td>
</tr>
<tr>
<td>RPC-TE</td>
<td>R</td>
<td>V</td>
<td>R</td>
</tr>
<tr>
<td>PC-TP</td>
<td>R</td>
<td>V</td>
<td>R</td>
</tr>
<tr>
<td>PC-TE</td>
<td>R</td>
<td>V</td>
<td>R</td>
</tr>
<tr>
<td>PV-TE</td>
<td>R</td>
<td>R</td>
<td>R</td>
</tr>
<tr>
<td>PV-TP</td>
<td>R</td>
<td>R</td>
<td>R</td>
</tr>
<tr>
<td>AV</td>
<td>R</td>
<td>R</td>
<td>R</td>
</tr>
<tr>
<td>DV</td>
<td>R</td>
<td>R</td>
<td>R</td>
</tr>
</tbody>
</table>

Table 1: Comparison of control results for Llull and Copeland elections and for plurality-rule elections. R means resistance to a particular control type and V means vulnerability. Results regarding plurality come from Bartholdi, Tovey, and Trick [BTT92] and Hemaspaandra, Hemaspaandra, and Rothe [HHR05].

Note that both Llull and Copeland have a higher number of constructive resistances, by two, than even Plurality, which was before this paper the reigning champ. (Although the results regarding plurality in Table 1 regard the unique winner version of control, for all the table’s Llull and Copeland cases our results hold both in the cases of unique winners and of nonunique winners, thus allowing an apples-to-apples comparison to hold.)
Table 1 separates its results into two main groups. In the first group we have results regarding procedural control where the chair affects candidate structure. In the second group the chair affects the voters. All our resistance results regarding candidate control follow via reductions from vertex cover, and all our vulnerability results follow via greedy algorithms. Our resistance results for the case of control by modifying voter structure follow from reductions—typically employing our UV technique—from variants of the exact-cover-by-3-sets problem.

Below we include two proof sketches somewhat indicative of our approaches and techniques.

Let us quickly define a certain notation. Within every election we fix some arbitrary order over the candidates, and any occurrence of a subset $D$ of candidates in a preference list means “the candidates from set $D$ are listed with respect to that fixed order.” Occurrences of $\overline{D}$ mean the same except that the candidates from $D$ are listed in the reverse order.

**Lemma 4 (subcase of Theorem 3)** Copeland is resistant to destructive control via deleting voters.

**Proof.** Let X3C-odd be a special case of the exact-cover-by-3-sets problem. We are given a set $B = \{b_1, \ldots, b_{3k}\}$ and a family of sets $\mathcal{S} = \{S_1, \ldots, S_n\}$ such that each $S_i$ is a size-3 subset of $B$. It is guaranteed that $k$ is odd. The question is whether it is possible to pick $k$ members of $\mathcal{S}$ in such a way that their union is equal to $B$. It is easy to see that this version of the X3C problem is NP-complete. Our proof follows by a reduction from X3C-odd to Copeland-DCDV.

Let $(B, \mathcal{S})$ be an X3C-odd instance with $|B| = 3k$ and $k = 2q + 1$. We construct the following Copeland election: We set the candidate set $C = \{p, U, V\} \cup B$, where $U$ is the despised candidate whom we want to prevent from being a winner. For every set $S_i$ we introduce $k + 1$ pairs of voters of types (i) and (ii) below.

(i) \[ U > V > S_i > p > B - S_i \]
(ii)-1 \[ B - S_i > p > U > V > \overline{S_i} \]
(ii)-2 \[ B - \overline{S_i} > U > p > V > S_i \]
(ii)-3 \[ B - S_i > V > p > U > \overline{S_i} \]

Note that voters of type (ii) come in three varieties. We introduce type (ii) voters in such a way that globally there are exactly $q$ voters of the second variety, exactly $q$ voters of the third one, and all the remaining type (ii) voters are of the first variety. It is irrelevant in which pairs particular varieties of type (ii) voters end up.

We claim that $U$ can be prevented from being a winner of this election via deleting at most $k$ voters if and only if $(B, \mathcal{S})$ is in X3C-odd.

Before any voters are deleted, we have the following results of head-to-head contests between each pair of candidates. Both $U$ and $V$ defeat $p$ by $2q = k - 1$ votes. $U$ defeats every candidate other than $p$ by more than $k$ votes. $V$ defeats every candidate other than $p$ and $U$ by more than $k$ votes. For any $i$, $1 \leq i \leq 3k$, $p$ is tied with $b_i$. For any $i$ and $j$, 9
\[1 \leq i < j \leq 3k, \ b_i \text{ is tied with } b_j.\] Therefore, it holds that \(U\) has Copeland score \(3k + 2\), \(V\) has Copeland score \(3k + 1\), and \(p\) and all members of \(B\) have Copeland score 0.

Since \(U\) has more than \(k\) votes of advantage over each candidate other than \(p\), via deleting \(k\) voters \(U\) can lose at most one vote, and this only happens if \(p\) defeats \(U\) after deleting \(k\) voters. Hence, the highest possible Copeland score in our election is \(3k + 2\), to ensure that \(U\) is not a winner after deleting \(k\) voters we need to guarantee that some candidate other than \(U\) defeats everyone else in head-to-head contest. The only candidate that can possibly achieve this is \(p\).

Thus we need to show that \(p\) can defeat all other candidates in our election via deleting \(k\) voters if and only if \((B, S)\) is in X3C-odd. Before any deletions, \(p\) loses \(k - 1\) votes to both \(U\) and \(V\). Thus any deleting of at most \(k\) candidates that guarantees \(p\) defeating both \(U\) and \(V\) has to include deleting only voters \(v\) that rank both \(U\) and \(V\) above \(p\). The only such voters are voters of type (i).

Thus if \(U\) can be prevented from being a winner then there exists a set \(V'\), \(|V'| \leq k\), of type (i) voters such that deleting voters \(V'\) guarantees that \(p\) defeats every \(b_i, 1 \leq i \leq 3k\), by at least one vote. However, regarding \(B\) candidates, deleting a single type (i) voter associated with set \(S_i\) gives \(p\) a one-vote advantage over exactly the members of \(S_i\). Thus \(V'\) defines an exact cover by 3-sets over \(B\). On the other hand, if the X3C-odd instance is positive then \(p\) clearly can be made a unique winner by deleting \(k\) voters.

The above proof illustrates some interesting points. For example, even though the proof is matched with a theorem that speaks of destructive control, it in fact can be seen to simultaneously prove that Copeland is resistant to constructive control via deleting voters in the unique winner case. Also, since in this problem the chair is not at liberty to change the preference lists (or tables) of the voters, we automatically get the analogous results for the case of irrational voters.

The proof of Lemma 4 is one of the simplest applications of the UV technique. The idea is that we introduce dummy candidates \(U\) and \(V\) and set up voters in such a way that to make a particular candidate win, say \(p\), we need to delete (or otherwise change, as in the applications of the UV technique to, e.g., bribery) only the voters who rank \(p\) below both \(U\) and \(V\). This allows us to introduce whatever padding candidates we need as long as we guarantee that \(U\) and \(V\) do not both defeat \(p\) on the padding voters' preference lists.

The next lemma illustrates a different technique.

**Lemma 5 (subcase of Theorem 3)** Copeland is resistant to constructive control via adding a bounded number of candidates.

**Proof.** The theorem follows via a reduction from the vertex cover problem to Copeland-k-CCAC. Our input is an undirected graph \(G = (V, E)\) and a nonnegative integer \(k\). The question is if there exists a subset \(W\) of \(V\) such that \(|W| \leq k\) and each edge \(e \in E\) has at least one of its endpoints in \(W\). We will show how this question can be reduced to asking whether within a certain Copeland election candidate \(p\) can be made a winner by adding at most \(k\) candidates.
W.l.o.g., we assume that $V = \{v_1, \ldots, v_n\}$ and that $E = \{e_1, \ldots, e_m\}$ and $k \geq 1$. We construct the following Copeland-k-CCAC instance: $C = \{p\} \cup E$ is the set of candidates that definitely participate in the election and $D = V$ is the set of spoiler candidates, i.e., the ones that the chair can convince to participate.

We set up the voters’ preference lists over $C \cup D$ in such a way that: (a) each $e_i$ defeats $p$, (b) each $e_i$ defeats all $v_i$ that $e_i$ is not incident to in $G$, (c) $p$ defeats each $v_i$, and (d) all the other pairs of candidates are tied in head-to-head contests. It is easy to construct such an election using at most polynomially (in the input size) many voters. For any two candidates $c'$ and $c''$ and a set of remaining candidates $K$, we can just introduce two voters, one with preference list $c' > c'' > K$ and the other one with preference list $K > c' > c''$. These two voters give $c'$ two votes advantage over $c''$ without affecting the relations between any other pair of candidates.

Before adding any of the spoiler candidates each candidate $e_i$ has Copeland score 1 and $p$ has Copeland score 0. Adding a single candidate $v_i$ has the following effect: (a) $p$ gets one additional Copeland point, (b) each $e_j$ candidate that is not incident with $v_i$ gets one additional Copeland point, and (c) $v_i$ has Copeland score 0. Thus the only possibility that $p$ has no less Copeland points than any other candidate in the election is that while adding candidates $v_i$, each $e_j$ happened to be incident with at least one of the added vertex-candidates. However, this means that the added candidates constitute a vertex cover of $G$. Thus if $p$ can be made a winner of this election by adding at most $k$ candidates then $G$ has a vertex cover of size at most $k$. The converse is trivial and the theorem is proven.

Resistance to control is generally viewed as a desirable property in system design. However, suppose one is in the role of someone trying to solve resistant control problems. Is there any hope? In fact, we have obtained a broad range of efficient algorithms for resistant-in-general control problems for the case when the number of candidates or voters is bounded. For example, each of the 16 problems regarding control via voters (Llull or Copeland; constructive or destructive; AV or DV or PV-TE or PV-TP) is in FPT (is fixed-parameter tractable, i.e., is not merely in P but indeed the family of P algorithms have degrees that do not depend on the value of the fixed number of voters or candidates) when the number of candidates is bounded, and also when the number of voters is bounded, and this claim holds even under the succinct input model (in which the voters are input via (preference-list, binary-integer-giving-frequency-of-that-preference-list) pairs). We prove some of these results using Lenstra’s algorithm for bounded-variable-cardinality integer programming. On the other hand, for irrational voters, all of the resistant candidate control problems remain resistant even for two voters.

5 Conclusions

We have shown that from the computational point of view the election systems of Llull and Copeland are broadly resistant to bribery and procedural control, regardless of whether the voters must have rational preferences. It is rather charming that Llull’s 700-year-old system shows perfect resistance to bribery and more resistances to (constructive) control than any
natural system (even far more modern ones) other than Copeland is known to possess, and this is even more remarkable when one considers that Llull’s system was defined long before control of elections was even explicitly studied. Copeland voting matches Llull’s perfect resistance to bribery and Llull’s broad resistance to (constructive) control.

References


